

Phantom Attack: Evading System Call Monitoring Rex Guo, Ph.D. Junyuan Zeng, Ph.D.

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web app RCE on joe-box and executed a reverse shell











An Incident ... A Defender's View



rule: untrusted program reads /etc/shadow condition:

```
syscall == open(at)
```

```
and read permission
```

and filename == /etc/shadow

and program is not in allowlist



- Introduction to System Call Monitoring
- Open Source System Call Monitoring Projects
- TOCTOU Phantom v1 Attack
- Semantic confusion Phantom v2 Attack
- Takeaways

System Call Monitoring (1)



System Call Monitoring – Syscall Interception (1)

- tracepoint/raw_tracepoint
 - Kernel static hook
 - tracepoint vs raw_tracepoint
 - Linux Kernel provides raw tracepoints: *sys_enter* and *sys_exit*
 - trace_sys_enter(struct pt_regs *regs, long id)
 - trace_sys_exit(struct pt_regs *regs, long id)



• Low overhead but only static system call interceptions

System Call Monitoring – Syscall Interception (2)

- kprobe/kretprobe
 - Dynamic hook in the kernel



- kprobe vs kretprobe
- Dynamic but slow compared to tracepoints and need to know exactly how data is placed on the stack and register
- LD_PRELOAD: not working in all cases
- Ptrace: performance overhead is high

System Call Monitoring – Syscall Data Collection

- Tracing programs collect system call data, e.g., arguments
- Tracing programs can "attach" to different hooks. When the hooks fire, tracing programs are executed
 - tracepoints/raw_tracepoints
 - kprobe/kretprobe
- Tracing programs implementations
 - Linux native mechanisms: *ftrace*, *perf_events* etc.
 - Kernel modules
 - eBPF programs: allow the execution of user code in the kernel

Open Source Projects (as of 07/15/2021)

- Falco (created by Sysdig)
 - One of the two security and compliance projects in CNCF incubating projects
 - The only endpoint security monitoring project in CNCF incubating projects
 - 3.9K github stars
 - It consumes kernel events and enriches them with information from the cloud native stack (e.g. Linux, containers, etc.)
 - Falco supports both kernel module and eBPF programs for tracing program implementation
- Tracee (created by Aqua Security)
 - 1.1K github stars
 - A runtime security and forensics tool based on eBPF

Vulnerabilities

- Time-of-check time-of-use (TOCTOU)
 - Time-of-check: tracing programs collect system call data (e.g. arguments)
 - Time-of-use: system call data used by kernel is different from what tracing programs check
 - e.g. sys_openat(int dfd, const char __user * filename, int flags, umode_t mode)
 - Phantom v1 attack exploits TOCTOU
- Semantic confusion
 - Kernel interprets data differently from the tracing programs
 - *e.g.* symbolic link is interpreted differently by the kernel and tracing programs
 - Phantom v2 attack exploits semantic confusion
- Falco is vulnerable to both Phantom v1 and v2
- Tracee is vulnerable to Phantom v1

https://dl.packetstormsecurity.net/1005-advisories/khobe-earthquake.pdf













TOCTOU – Falco

- <u>CVE-2021-33505</u> CVSS v3.0 score 7.3
- Falco older than v0.29.1 (or open source sysdig)
 - Commercial versions based on the open source agent are also affected (confirmed by the open source maintainer)
- It uses *raw tracepoints* (*sys_enter* and *sys_exit*) to intercept syscalls
- User space pointers are read directly by its tracing programs
 - In the implementations of both kernel module and eBPF programs

TOCTOU – Falco

• We evaluated the important syscalls in Falco rules.

Syscall	Category	TOCTOU?	
connect	Network	Y	
sendto/sendmsg	Network	Y	
open/openat	File	Y	
execve	File	Ν	
rename	File	Y	
renameat/renameat2	File	Y	
mkdir	File	Y	
mkdirat	File	Y	
rmdir	File	Y	
unlink/unlinkat	File	Y	
symlink/symlinkat	File	Y	
chmod/fchmod/fchmodat	File	Y	
creat	File	Ŷ	

TOCTOU – Tracee

- Tracee (v0.4.0) is vulnerable to TOCTOU for many system calls, e.g., connect syscall, etc.
- No CVE given. Here are some quotes from the maintainers:
 - "As you probably know, TOCTTOU attacks on system calls wrappers(/tracers) is a well known issue, and Tracee is no exception."
 - "And yes we agree on the fact that there's no CVE or novel finding and therefore you could talk about it publicly."
 - Interpret yourself 😳

Phantom v1 Exploit Plan (Sys_exit is Monitored)

- Triggers the target system call with malicious arguments
- Let kernel reads the malicious arguments and performs the intended malicious action
- Overwrites the data structure pointed by the user space argument pointer with benign data
- At sys_exit, tracing program reads the data structure pointed by the user space pointer and checks against the rules
- Challenges:
 - When does the kernel thread reads it?
 - How can we synchronize the overwrite with the kernel thread read?
 - Are the racing windows big enough for each syscalls?
 - How to ensure the tracing program get the overwritten copy?

Userfaultfd Syscall

- Normally page faults are a kernel internal thing...
 - Why offload page faults to userland?
- Memory externalization: running programs with memory residing on a remote node
 - Memory is transferred from the memory node to the compute node on access



 Memory can be transferred from the compute node to the memory node if it's not frequently used during memory pressure



• Once userfaultfd triggers, kernel thread is paused and waits for user space response

Helps exploitation on kernel race condition bugs

Interrupts and Scheduling

- An interrupt notifies the processor with an event that requires immediate attention
- An interrupt diverts the program control flow to an interrupt handler
- Interrupt can be triggered indirectly from system calls
 - Hardware interrupts (networking, e.g., connect)
 - Interprocessor interrupts (IPIs) (e.g., mprotect)
- sched_setscheduler()
 - set SCHED_NORMAL / SCHED_IDLE
 - Realtime policies require CAP_SYS_NICE or Realtimekit]
- sched_setaffinity(): pin task to CPU bitmask



userfaultfd thread

any CPU

overwrite thread

pin CPU 2

block on cond mutex

main thread

pin CPU 3 mmap page A register userfaultfd thread

> sys_enter openat *filename -> page A

kernel thread triggers page fault

*filename -> page A (not allocated)

	main thread	userfaultfd thread	overwrite thread
pin CPU 3 mmap page A register userfaultfd thread	any CPU	pin CPU 2	
		block on cond mutex	
*filename -> page A (not allocated)	sys_enter openat *filename -> page A		
*filename -> malicious filename	kernel thread triggers page fault	write page A with malicious filename	
		release cond mutex	
		IOCTL return execution back to kernel	write benign name
,	7		





Semantic Confusion – File Link

- Semantic confusion indicates the different interpretations of data (e.g. system call arguments) between the kernel and tracing programs
- In particular for file link interpretation, kernel trails symbolic link to actual file while tracing programs read the link without any interpretations
- Falco is vulnerable to semantic confusion
 - It reads symbolic link without any interpretation
 - No CVE because symlink(at) and link(at) are monitored
 - But practically the detection team need to track all symlink(at)/link(at) to any file based rules ⁽³⁾
- Tracee is not vulnerable to openat
 - *security_file_open* LSM hook: *filename* has been interpreted by the kernel

Phantom v2 – An File Link Example

```
rule: untrusted program reads /etc/shadow
condition:
    syscall == open(at)
    and read permission is used
    and filename == /etc/shadow
    and not program is not in allowlist
```

- Steps to bypass the Falco rule
 - Create a symlink /tmp/shadow -> /etc/shadow
 - Tracing programs read the symlink /tmp/shadow
 - Syscall openat monitoring reports /tmp/shadow is opened
 - Rule is bypassed

Mitigation

- Detection (Falco team)
 - Detect (unprivileged) usage of the `userfaultfd` syscall (Implemented)
 - Detect a user registering a memory address range
 - Detect a user copying a continuous memory chunk into the userfaultfd registered range and (optionally) wake up the blocked thread (kernel)
- Read the data used by system calls
 - LSM hook: a list of check points

LSM hook	used by Tracee v0.4.0	Protected syscall
security_b	orm_check	execve, execveat
security_fil	e_open	open, openat
security_in	ode_unlink	unlink, unlinkat
security_m	map_addr	mmap, mmap_pgoff
security_fi	e_mprotect	mprotect

• Kernel data structure: read arguments of *execve* from *mm->arg_start*

Takeaways

- Phantom attack is generic and exploits the fact that kernel and tracing programs
 - Can read data at different times (Phantom v1)
 - Can interpret data differently (Phantom v2)
- Kernel raw tracepoints on system calls are not ideal for secure tracing
- Other tracing implementations can be vulnerable. E.g., kprobe
- Mitigation:
 - Detect abnormal usages of userfaultfd
 - Ensure kernel and secure tracing programs (1) Read the same data (2) Interpret data in the same way
- If you are interested in discussing further:
 - @Xiaofei_REX (OpenDM)
 - https://github.com/rexguowork/phantom-attack (will be released during Defcon)

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